## Implementation of ALGOL 60 for the English Electric KDF9

By F. G. Duncan

Our decision to implement ALGOL 60 for the KDF9 was taken about 18 months ago. We saw in ALGOL 60 the possibility of its use as an automatic programming language for a wide range of users, particularly those who had become accustomed to working with earlier languages. On the one hand, ALGOL 60 was quite obviously much more powerful than any of the others, and on the other hand we wanted to make a contribution towards the effort to establish a language that could be universally understood and implemented. From the first our ambition was to implement the ALGOL 60 Report (Naur, ed., 1960) as fully as we possibly could. Experience in writing and using an "optimizing" compiler for DEUCE (Duncan and Hawkins, 1959; Duncan and Huxtable, 1960) led to the desire to combine full ALGOL translation with the generation of "efficient" (or relatively fast-operating) object programs. We realized, of course, that this would entail considerable original work, and planned to put about five man-years into the project.

This was at Kidsgrove, in the Data Processing and Control Systems Division. There was also an interest in ALGOL in the Atomic Power Division at Whetstone, where, in fact, there were people who had produced real, working, compilers. They were prepared to put in one man-year or so on an independent ALGOL compiler, to be made available before the more elaborate Kidsgrove version. This would both help us to get ALGOL established among users, and provide checks on the development of the elaborate scheme. At first we saw it as a very restricted ALGOL compiler; but following some informal contacts with Professor van Wijngaarden and Dr. Dijkstra of Amsterdam, whose compiler (Dijkstra, 1961)\* had by this time been completed, the Whetstone people were entertained magnificently at Amsterdam for a week, and returned able to remove most of the restrictions they had previously imposed.

This occasion is a suitable one for us to express publicly our thanks to the Mathematisch Centrum, Amsterdam, for their great generosity in making sure that their work was fully understood by us; and also to Dr. Naur's group at Copenhagen, with whom I spent a very useful week rather more than a year ago. Our relations with both these groups have always been very friendly.

Towards the end of last year, therefore, we found ourselves with the possibility of having two compilers, both able to deal with almost the whole generality of the ALGOL Report. Naturally the question arose as to whether one or other of these projects should be

\* See also the paper on p. 125 of this issue.—Ed.

dropped. We decided to continue with both. Our earlier experience and our discussions with a variety of prospective users indicated that there would be a place for each of two compilers with complementary characteristics.

## **Compiler Characteristics**

These characteristics, briefly, are as follows:

Whetstone. The aim is fast compiling. The operation is what is sometimes called "one-pass-load-and-go." There is no particular attempt to obtain efficiency in the object program. As in the case of the Elliott compiler described by Mr. Hoare,\* compilation time is much the same as paper-tape reading time. It is thought that there will be about 3,000 words of instructions in the compiler, and we hope that it will be available during the autumn of 1962.

Kidsgrove. Compiling takes longer, but the time should not be more than three minutes. There are several passes through the program, aimed at recognizing and giving special treatment to certain situations which are amenable to "optimization." Examples of the situations are simple "for" statements making simple use of the controlled variable, and exceptionally well-behaved procedure declarations and their corresponding calls. Situations which cannot clearly be recognized as qualifying for optimization are given the full general treatment—a "fail safe" method. An account of the detection processes is given by Hawkins and Huxtable (1962).

Now for the properties which the two compilers have in common. First, both accept identical versions of ALGOL 60. The restrictions are:

- 1. No integer labels.
- 2. A complete specification part must be provided with each procedure declaration.
- 3. The declarator own is given the sense of Reformulation 23 of *ALGOL Bulletin* 14, except that we do not allow own arrays with "dynamic" bounds.
- 4. Where a formal parameter is specified as a procedure, all the corresponding actual procedures must have identical specifications. Second, both will be able to accept the same paper-tape versions of the ALGOL texts. Third, both will be able to deal with the same procedures whose bodies are in KDF9 User Code. This is an assembly code for the machine whose instructions are in one-to-one correspondence with those of the computer itself. This point will be expanded later.

Thus there should be complete two-way compatibility. Here one might mention an experimental compiler for

<sup>\*</sup> See p. 127.

DEUCE (a wonderful old machine, which still provides solace and recreation for those of us who like to be reminded occasionally of what real programming really was). This compiler has just been provided with routines for dealing with KDF9 paper tapes, so that we have a one-way compatibility between two entirely different machines—programs for DEUCE ALGOL will be able to run without change on KDF9.

As for size of machine, both compilers require a core store of 8,192 words. Whetstone would like three magnetic-tape units and Kidsgrove four, although both systems, we think, can be operated with two, with some inconvenience and fewer "facilities."

There are some important features of the Kidsgrove scheme which we do not anticipate will be available with the Whetstone compiler. They are embodied in an operating system which is still under development. This provides for several modes of operation—compileand-run, compile several programs, compile mixed sequence of ALGOL and User Code programs, and so on. It will deal with amendments to programs, and the printing out of "useful information." The input stage to the compiler permits the ALGOL text to be presented in almost any hardware representation within reason. The user is only required to "declare" his code—that is, provide a list of the actual tape characters, or combinations of characters, corresponding to the basic symbols of the Report. A testing system is envisaged, the principle being that one should communicate with the machine only in ALGOL. Calls of special test procedures could be included in a "test" compilation and ignored in a "run" compilation.

We hope that the Kidsgrove scheme, or at least the compiler, will be working by the end of 1962. It seems there will be around 20,000 words of instructions. Both compilers are being written in User Code.

We intend to provide a library of ALGOL procedures to go with either compiler. Naturally enough, the Kidsgrove scheme will provide the more elaborate facilities for handling this library. Library procedures will have their bodies either in ALGOL or in User Code. The content of the library will depend, of course, on contributions from users, and on what can be taken from periodicals or other publications; we would try to provide a reasonable start with standard functions and input and output routines.

### Use and Teaching

As to the use of the two compilers, I can see the Whetstone compiler being popular in places where jobs are "one-off"—where the requirement is to get on and off the machine quickly and at very short notice—where, for example, methods are under investigation. The Kidsgrove scheme is more suitable for places where operating schedules can be planned and where programs are run for "production" purposes, where efficient use of computer time is regarded as important. Extreme cases of both situations can be found among our customers.

Some effort has been put into teaching ALGOL at both Kidsgrove and Whetstone. The teaching manual (Green, 1961) which was produced at Kidsgrove has been well received outside English Electric, and most of it could be useful to those who would like a gentle conducted tour of the essential features of the language. Perhaps it is too gentle and too cautious. I have found difficulty in making an ALGOL course, including recursive procedures, Jensen's device, and side effects, last as long as four days. I believe that experience at other establishments has been similar. We have a few self-taught ALGOL programmers, one of whom has produced a program for structural analysis running to a fair-sized book. It makes use of recursive procedures, naturally.

#### Non-ALGOL Procedures

Earlier in this talk procedures with bodies in User Code were mentioned. I think our work on this topic is worth describing; it could be adapted for use with other machines.

Section 5.4.6 of the ALGOL Report provides that procedure bodies may be expressed in non-ALGOL code. In our case we have taken KDF9 User Code, with some extensions to provide communication with the surrounding ALGOL context. Such communication is solely through the parameter list. Since the scheme is to be implemented for two entirely different compilers, code procedures must not make use of any internal properties of a compiler or of the object program.

We make the required extension to the syntax of section 5.4.6 as follows:

<description> :: = (no. of nesting cells used>, (no. of
 Q-stores used>, (no. of W-stores used>, (no. of
 V-stores used>

(instruction) :: = (KDF9 User Code instruction) |
 "(formal parameter)" | = "(formal parameter)" |
 J''(formal parameter)"

<instruction sequence> :: = <instruction> | <instruction
sequence>; <instruction>

⟨code⟩ :: = KDF9 ⟨description⟩; ALGOL

Semantically, specifiers are restricted to exclude switches and procedures. The three forms for instruction beyond those in ordinary user code are respectively for "fetch," "store," and "jump"—the formal parameter for the last must be a label. Although the example below—a User-code version of the Innerproduct procedure given on p. 34 of the ALGOL report—can be completely understood only by those familiar with KDF9, it should convey the essentials of the scheme to anyone who understands ALGOL.

## Example

procedure Innerproduct (a, b) Order: (k, p) Result: (y);
value k;

```
integer k, p; real y, a, b;
KDF9  4, 0, 0, 2;
        ZERO; DUP; DUP; =V1; =V2;
2; SET 1; +; ="p";
        "p"; "k"; -; J1 > Z;
        V2; V1; "a"; "b"; ×+F; =V1; =V2;
        "p"; J2;
1; V2; V1; ROUND F; ="y";
```

### ALGOL

The formal parameters, which are each enclosed in quotation marks in the KDF9 text, are replaced automatically by sequences of instructions for getting access to the required quantities. The system allows calls by name and by value. In the example, V1 and V2 correspond to the local variable s of the Report version. In this version the sum is accumulated double length and then rounded.

The reason for providing so fully for machine-code procedures is to simplify the introduction of features which it would otherwise be inconvenient or impossible to express within ALGOL. For example, one might need a set of procedures for evaluating some frequentlyused functions, and the running speed of the ALGOL versions, even when translated by an optimizing compiler, might not be sufficiently close to that of the corresponding machine-code versions. Input and output, and magnetictape procedures must either themselves be in machine code or make use of machine-code procedures. Something like the scheme we have proposed is necessary if one wishes to get beyond the stage of "read one number, punch one number."

As was mentioned earlier, machine-code procedures and ALGOL procedures can be included in the library. It follows that the user of library procedures for input and output need know nothing about the KDF9 User Code.

A matrix scheme proposed for use with our ALGOL system (Denison, 1962) makes use of a number of procedures which have already been expressed in ALGOL. It is probable that they will be rewritten in User Code for the sake of speed.

## Acknowledgement

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# **Operating experience with FORTRAN**

## By A. E. Glennie

My purpose in this talk will be to describe the lessons that I have learned from my own experience, and that of my colleagues, in using FORTRAN during the last three years. Some of the points I shall make are specifically about the FORTRAN language itself; others are about automatic coding in general, and computer operating systems incorporating compilers. I hope that what I have to say about the latter aspects, as revealed in the use of FORTRAN, may be of interest and value to those of you whose interests and preferences may be in other languages—or in fields other than scientific computation.

### The History of Our Use of FORTRAN

I think that you will be interested to hear how the use of FORTRAN in the U.K.A.E.A. developed, as our story is quite typical of the evolution of a laboratory's technique and practice. When, in 1958, we started our first experiments with FORTRAN on the IBM 704, we had had a long tradition of machine-language coding, but had also had some experience of automatic coding.

We were not, I think, prejudiced against automatic coding, yet we were somewhat disappointed with our first experience with FORTRAN. This was the